

Customer Notification

IE-703288-G1-EM1TM

In-Circuit-Emulator

Operating Precautions

Target Device

V850ES/SG2

V850ES/SJ2

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(A) Table of Operating Precautions

No.	Outline	IE-703288-G1-EM1			
		Rev.	1.02	1.03	1.04
		Control-Code ^{Note}	C	D	E
1	ROM correction function cannot be emulated (Direction of use)		X	X	X
2	Use-prohibited area (Direction of use)		X	X	X
3	Emulator memory settings (Direction of use)		X	X	X
4	Accuracy of ADC and DAC (Technical limitation)		X	✓	✓
5	Watchdog timer during break (Direction of use)		X	X	X
6	Timer M during break (Direction of use)		X	X	X
7	Timer M compare interrupt (Specification change notice)		X	X	X
8	Access of UAnRX register during break (Specification change notice)		X	X	X
9	Access of CBnRX register during break (Specification change notice)		X	X	X
10	Access of C0RGPT register during break (Specification change notice)		X	X	X
11	Access of C0TGPT register during break (Specification change notice)		X	X	X
12	Access of C0GNCTRL register during break (Specification change notice)		X	X	X
13	SLD instruction precaution (Specification change notice)		X	X	X
14	aFCAN transmission / reception (Technical limitation)		X	X	✓
15	TMPn / TMQn external event counter function (Direction of use)		X	X	X
16	TMPn / TMQn capture operation (Direction of use)		X	X	X

Operating Precautions for IE-703288-G1-EM1

No.	Outline	IE-703288-G1-EM1			
		Rev.	1.02	1.03	1.04
		Control-Code ^{Note}	C	D	E
17	aFCAN: Rx limitation (Technical limitation)		X	X	X

✓ Not applicable

X Applicable

Note: The Control Code is indicated by the letter appearing at the 2nd position from the left in the serial number of the product.

(B) Description of Operating Precautions

No. 1	ROM correction function cannot be emulated (Direction of use)
	<u>Details</u> The ROM correction function cannot be emulated. <u>Workaround</u> There is no workaround.

No. 2	Restriction on use-prohibited area (Direction of use)				
	<u>Details</u> A fail-safe break is not generated if program execution or an access to a use prohibited memory area of the device is attempted.				
	<u>Workaround</u> A break can be generated when the program is executed or a memory access occurs by setting a break in the debugger under the following conditions:				
	Detailed description of restriction on use-prohibited area				
	Device:	D703260(Y), D703270(Y), D703280(Y)	D70(F)3261(Y), D70(F)3271(Y), D70(F)3281(Y), D70(F)3264(Y), D70(F)3274(Y), D70(F)3284(Y)	D703262(Y), D703272(Y), D703282(Y), D703265(Y), D703275(Y), D703285(Y), D703287(Y)	D70(F)3263(Y), D70(F)3273(Y), D70(F)3283(Y), D70(F)3266(Y), D70(F)3276(Y), D70(F)3286(Y), D70(F)3288(Y)
	Use-prohibited area in which a fail safe break does not occur:	(1): 0x3FF8000 to 0x3FF8FFF	(2): 0x60000 to 0x7FFFF (3): 0x3FF0000 to 0x3FF6FFF	(4):0x3FF0000 to 0x3FF4FFF	(5):0xA0000 to 0xFFFFF (6):0x3FF0000 to 0x3FF2FFF
Execution access break (used to replace fail save break):	<ul style="list-style-type: none">Event: ExecutionAddress: Area (1) (2 execution events used for above conditions)	<ul style="list-style-type: none">Event: ExecutionAddress:Area (2)Event: ExecutionAddress:Area (3) (4 execution events used for above conditions)	<ul style="list-style-type: none">Event: ExecutionAddress: Area (4) (2 execution events used for above conditions)	<ul style="list-style-type: none">Event: ExecutionAddress:Area (5)Event: ExecutionAddress:Area (6) (4 execution events used for above conditions)	
R/W access break (used to replace fail save break):	<ul style="list-style-type: none">Event: R/WAccess size: No ConditionAddress: Area (1) (2 execution events used for above conditions)	<ul style="list-style-type: none">Event: R/WAccess size: No ConditionAddress: Area (3) (2 access events used for above conditions)	<ul style="list-style-type: none">Event: R/WAccess size: No ConditionAddress: Area (4) (2 access events used for above conditions)	<ul style="list-style-type: none">Event: R/WAccess size: No ConditionAddress: Area (6) (2 access events used for above conditions)	

No. 3	Emulator memory settings (Direction of use)
	<p><u>Details</u></p> <p>When the debugger connects to the emulator the memory settings of the devicefile are not automatically set for several devices. This applies for devices with internal memory sizes that do not match to the below list.</p> <p>The size of internal memory of the emulator can only be set to the following values:</p> <p>Internal ROM: 32 KBytes, 64 KBytes, 128 KBytes, 256 KBytes, 512 KBytes or 1 MByte.</p> <p>Internal RAM: 4 KBytes, 12 KBytes, 28 KBytes or 60 KBytes.</p> <p><u>Greenhills Multi:</u></p> <p>For devices which have different memory sizes than listed above, the size of internal ROM or RAM memory is set to the next smaller size of the above list.</p> <p>E. g.: for devices that contain 384/32 KBytes of internal ROM/RAM memory, 256/28 KBytes is set (512/28 KBytes is set for devices containing 640/48 KBytes of internal ROM/RAM).</p> <p><u>IAR Embedded workbench:</u></p> <p>When connecting the debugger to the emulator an error message is shown if the ROM size selected in the devicefile does not fit the possible settings of the emulator (see list above). The RAM size is automatically set to the next larger acceptable option.</p> <p><u>Workaround</u></p> <p>Change the settings when the connection between debugger and emulator is established.</p> <p><u>Greenhills Multi:</u></p> <p>Set the desired size of internal ROM/RAM memory in the .rc file which is executed when the debugger connects to the emulator using the target command "CPU [R= A=]".</p> <p>Set the internal memory to the next larger size of the above list (e. g. select 512 KBytes for a devices that contains 384 KBytes of internal ROM).</p> <p>The actual settings for internal memory can be checked in the target window using the "CPU" command.</p> <p><u>IAR Embedded workbench:</u></p> <p>Set the internal memory size in the configuration window for "Hardware Settings" which opens automatically when connecting to the emulator for the first time. The "Hardware Settings" window can also be invoked under the "Emulator Hardware Setup" drop down menu to alter or check the actual settings. Set the internal memory to the next larger size of the above list (e. g. select 512 KBytes for a devices that contains 384 KBytes of internal ROM).</p> <p>Take care that the memory borders of the actual device are not exceeded since the actual memory size provided by the emulator may be larger than the memory size of the emulated device.</p>

No. 4	Accuracy of A/D converter and D/A converter (Technical limitation)
	<p><u>Details</u></p> <p>The accuracy of the A/D converter does not meet the specification.</p> <p><u>Workaround</u></p> <p>There is no workaround. The accuracy has been improved with emulation boards of control code D or later (the conventional error is approx. 8%).</p>

No. 5	Watchdog timer during break (Direction of use)
	<p><u>Details</u></p> <p>When both of the following conditions (a) and (b) are fulfilled simultaneously and a break occurs, the watchdog timer does not stop and will cause a reset or non maskable interrupt. If a reset occurs, the debugger hangs up.</p> <p>Conditions that need to be fulfilled so that the above behaviour occurs:</p> <p>(a) The main clock or subclock is selected as the clock source of the watchdog timer and</p> <p>(b) The ring oscillator is stopped (RSTOP flag = 1).</p> <p><u>Workaround</u></p> <p>As a workaround to prevent the above behaviour do not stop the ring oscillator clock.</p>

No. 6	Timer M during break (Direction of use)
	<p><u>Details</u></p> <p>When a break occurs while the following conditions (a) and (b) are both fulfilled, timer M does not stop even if the peripheral break function has been set to 'break'.</p> <p>(a) INTWT, Ring oscillator clock (fR/8) or subclock is selected as the clock source for timer M.</p> <p>(b) The main clock is stopped by setting the MCK flag.</p> <p>(Note: The peripheral break function is not supported by the debugger ID850 V2.51.)</p> <p><u>Workaround</u></p> <p>Implement one of the below workarounds to stop timer M during a break using the peripheral break function:</p> <p>(a) Use the main clock (fXX, fXX/2, fXX/4, fXX/64, fXX/512) as the source clock for timer M.</p> <p>(b) Do not stop the main clock oscillation.</p>

No. 7	Timer M compare interrupt (Specification change notice)
	<p><u>Details</u></p> <p>An unexpected interrupt occurs after activation of timer M when the compare register TM0CMP0 contains the value 0xFFFF.</p> <p>The diagram illustrates the timing of the Timer M compare interrupt. It shows four signals: TM0CE, TMM (16bit counter), TM0CMP0, and INTTMEQ0. TM0CE is a pulse that starts the timer. TMM is a 16-bit counter that starts at 0000H and increments. TM0CMP0 is a register that is set to 0xFFFF. INTTMEQ0 is the interrupt output. An arrow labeled 'TMM starts to operate' points to the start of the TMM counter. The TMM counter values are shown as 0000H, 0001H, 0002H, ..., FFFE, FFFF, 0000H. The TM0CMP0 register is set to 0xFFFF. The INTTMEQ0 signal shows a pulse that occurs when the TMM counter reaches 0000H, which is unexpected because the compare register is 0xFFFF. This pulse is circled and labeled 'Interrupt'.</p> <p><u>Workaround</u></p> <p>Do not set TM0CMP0 to 0xFFFF.</p>

No. 8	Access of UAnRX register during break (Specification change notice)
	<p><u>Details</u></p> <p>An overrun error occurs under the following conditions (a) to (c):</p> <p>(a) If a break occurs after reading the UART receive buffer register (UAnRX) and the UAnRX register is displayed in the I/O register window of the debugger, an overrun error occurs when UART reception is performed for the next time.</p> <p>(b) If a software break occurs immediately after reading the UART receive buffer register (UAnRX), an overrun error occurs when UART reception is performed the next time regardless of whether or not the UAnRX register is displayed in the I/O register window.</p> <p>(c) If a DMA transfer from the UART receive buffer register (UAnRX) is performed during a break NOTE, an overrun error occurs when UART reception is performed the next time.</p> <p>Note: Including breaks by the RAM monitor function or DMM function. However the realtime RAM monitor function does not cause this behaviour since it does not set breaks.</p> <p>Remark: An overrun error also occurs when the UART receives data multiple times during a break (This complies with the specification of the emulator).</p> <p><u>Workaround</u></p> <p>(a) Do not display the UAnRX register in the I/O register window.</p> <p>(b) Set a hardware break when setting a break immediately after reading the UAnRX register</p> <p>(c) There is no workaround.</p>

No. 9	Access of CBnRX register during break (Specification change notice)
	<p><u>Details</u></p> <p>When the CSIBn receive data register (CBnRX) is read, it usually starts the next reception operation. Under the following conditions (a) and (b), however, the next reception operation is not started even if CBnRX is read.</p> <p>(a) If a software break occurs immediately after reading the CSIBn receive register (CBnRX).</p> <p>(b) If a DMA transfer from the CSIBn receive data register (CBnRX) is performed during a break NOTE. As a result the communication stops or the DMA controller stops.</p> <p>Note: Including breaks by the RAM monitor function or DMM function. However the realtime RAM monitor function does not cause this behaviour since it does not set breaks.</p> <p><u>Workaround</u></p> <p>(a) Set a hardware break when setting a break immediately after reading the CBnRX register.</p> <p>(b) There is no workaround.</p>

No. 10	Access of C0RGPT register during break (Specification change notice)
	<p><u>Details</u></p> <p>Under the following conditions (a) and (b), the read pointer (RGPT) that should be incremented is not incremented and the same data as previously read is read again.</p> <p>(a) If a software break occurs immediately after reading the CAN0 module receive history list register (C0RGPT)</p> <p>(b) If a DMA transfer from the CAN0 module receive history list register (C0RGPT) is performed during a break NOTE.</p> <p>Note: Including breaks by the RAM monitor function or DMM function. However the realtime RAM monitor function does not cause this behaviour since it does not set breaks.</p> <p><u>Workaround</u></p> <p>(a) Set a hardware break when setting a break immediately after reading the C0RGPT register.</p> <p>(b) There is no workaround.</p>

No. 11	Access of C0TGPT register during break (Specification change notice)
	<p><u>Details</u></p> <p>Under the following conditions (a) and (b), the read pointer (TGPT) that should be incremented is not incremented and the same data as previously transmitted is transmitted again.</p> <p>(a) If a software break occurs immediately after reading the CAN0 module transmit history list register (C0TGPT).</p> <p>(b) If a DMA transfer from the CAN0 module transmit history list register (C0TGPT) is performed during a break NOTE.</p> <p>Note: Including breaks by the RAM monitor function or DMM function. However the realtime RAM monitor function does not cause this behaviour since it does not set breaks.</p> <p><u>Workaround</u></p> <p>(a) Set a hardware break when setting a break immediately after reading the C0TGPT register.</p> <p>(b) There is no workaround.</p>

No. 12	Access of C0GNCTRL register during a break (Specification change notice)
	<p><u>Details</u></p> <p>When a register access is performed in the following sequence, an unexpected forcible shutdown may occur after the sequence is complete.</p> <p>Sequence :</p> <ol style="list-style-type: none"> (1) The EFSD bit of the CAN0 module control register (C0GMCTRL) is set. (2) The I/O register^{NOTE} is accessed. (3) The GOM bit of the CAN0 mode control register (C0GMCTRL) is cleared. <p>Note: I/O register access except for clearing the GOM bit of the C0GMCTRL register</p> <p>The conditions under which a forcible shutdown takes place are shown below:</p> <ol style="list-style-type: none"> (a) If a break occurs immediately after the I/O register access in (2) occurs. (b) If a break by the RAM monitor function or the DMM function occurs immediately after the I/O register access in (2) occurs. (c) Stepwise execution is performed for the I/O register access in (2). <p><u>Workaround</u></p> <p>Be sure to set the EFSD bit and clear the GOM bit successively when executing a forcible shutdown. Do not perform a register access in the above sequence when not performing a forcible shutdown.</p>

No. 13	SLD instruction precaution (Specification change notice) (Specification change notice)																																				
	<p><u>Details</u></p> <p>If a conflict occurs between the decode operation of the instruction (<2> in the examples mentioned below) immediately before the sld instruction (<3> in the examples) following a special instruction (<1> in the examples) and an interrupt request before execution of the special instruction is complete, the execution result of the special instruction may not be stored in a register as expected.</p> <p>This situation may only occur when the same register is used as the destination register of the special instruction and the sld instruction, and when the register value is referenced by the instruction followed by the sld instruction.</p> <p><u>Conditions under which the conflict occurs:</u></p> <p>The situation may occur when all the following conditions (1) to (3) are satisfied.</p> <p>(1) Either condition (I) or (II) is satisfied</p> <p>Condition (I): The same register is used as the destination register of a special instruction (see below) and the subsequent sld instruction and as the source register (reg1) of an instruction shown below followed by the sld instruction (See Example 1).</p> <table><tr><td>mov reg1,reg2</td><td>not reg1,reg2</td><td>satsubr reg1,reg2</td><td>satsub reg1,reg2</td></tr><tr><td>satadd reg1,reg2</td><td>or reg1,reg2</td><td>xor reg1,reg2</td><td>and reg1,reg2</td></tr><tr><td>tst reg1,reg2</td><td>subr reg1,reg2</td><td>sub reg1,reg2</td><td>add reg1,reg2</td></tr><tr><td>cmp reg1,reg2</td><td>mulh reg1,reg2</td><td></td><td></td></tr></table> <p>Condition (II): The same register is used as the destination register of a special instruction (see below) and the subsequent sld instruction and as the source register (reg2) of an instruction shown below followed by the sld instruction (See Examples 2 and 3).</p> <table><tr><td>not reg1,reg2</td><td>satsubr reg1,reg2</td><td>satsub reg1,reg2</td><td>satadd reg1,reg2</td></tr><tr><td>satadd imm5,reg2</td><td>or reg1,reg2</td><td>xor reg1,reg2</td><td>and reg1,reg2</td></tr><tr><td>tst reg1,reg2</td><td>subr reg1,reg2</td><td>sub reg1,reg2</td><td>add reg1,reg2</td></tr><tr><td>add imm5,reg2</td><td>cmp reg1,reg2</td><td>cmp imm5,reg2</td><td>shr imm5,reg2</td></tr><tr><td>sar imm5,reg2</td><td>shl imm5,reg2</td><td></td><td></td></tr></table> <p>Special instruction:</p> <ul style="list-style-type: none">• ld instruction: ld.b, ld.h, ld.w, ld.bu, ld.hu• sld instruction: sld.b, sld.h, sld.w, sld.bu, sld.hu• Multiply instruction: mul, mulh, mulhi, mulu <p>(2) When the execution result of the special instruction (see above) has not been stored in the destination register before execution of the instruction (instruction of condition (I) or (II)) immediately before the sld instruction starts in the CPU pipeline.</p>	mov reg1 ,reg2	not reg1 ,reg2	satsubr reg1 ,reg2	satsub reg1 ,reg2	satadd reg1 ,reg2	or reg1 ,reg2	xor reg1 ,reg2	and reg1 ,reg2	tst reg1 ,reg2	subr reg1 ,reg2	sub reg1 ,reg2	add reg1 ,reg2	cmp reg1 ,reg2	mulh reg1 ,reg2			not reg1, reg2	satsubr reg1, reg2	satsub reg1, reg2	satadd reg1, reg2	satadd imm5, reg2	or reg1, reg2	xor reg1, reg2	and reg1, reg2	tst reg1, reg2	subr reg1, reg2	sub reg1, reg2	add reg1, reg2	add imm5, reg2	cmp reg1, reg2	cmp imm5, reg2	shr imm5, reg2	sar imm5, reg2	shl imm5, reg2		
mov reg1 ,reg2	not reg1 ,reg2	satsubr reg1 ,reg2	satsub reg1 ,reg2																																		
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No. 13	SLD instruction precaution (Specification change notice) (Specification change notice)						
	<p>(cont.)</p> <p>(3) When the decode operation of the instruction (instruction of condition (I) or (II)) immediately before the sld instruction and interrupt request servicing conflict.</p> <p><u>Examples of instruction sequences that may cause the conflict:</u></p> <p>Example 1:</p> <table border="0"> <tr> <td style="vertical-align: top;"> <pre> <1> ld.w [r11], r10 : <2> mov r10, r28 <3> sld.w 0x28, r10 </pre> </td><td style="vertical-align: top;"> <p>This situation occurs when the decode operation of the mov instruction (<2>) immediately before the sld instruction (<3>) and interrupt request servicing conflict before the execution of the special instruction ld (<1>) is complete.</p> </td></tr> </table> <p>Example 2:</p> <table border="0"> <tr> <td style="vertical-align: top;"> <pre> <1> ld.w [r11], r10 : <2> cmp imm5, r10 <3> sld.w 0x28, r10 <4> bz label </pre> </td><td style="vertical-align: top;"> <p>This situation occurs when the decode operation of cmp (<2>) immediately before the sld instruction (<3>) and interrupt request servicing conflict before execution of the special instruction ld (<1>) is complete. As a result, the compare result of the cmp instruction becomes undefined, which may cause an unexpected operation of the branch instruction bz (<4>).</p> </td></tr> </table> <p>Example 3:</p> <table border="0"> <tr> <td style="vertical-align: top;"> <pre> <1> ld.w [r11], r10 : <2> add imm5, r10 <3> sld.w 0x28, r10 <4> setf c, r16 </pre> </td><td style="vertical-align: top;"> <p>This situation occurs when the decode operation of the add instruction (<2>) immediately before the sld instruction (<3>) and interrupt request servicing conflict before execution of the special instruction ld (<1>) is complete. As a result, the result of the add instruction and the depending status flags become undefined, which may cause an unexpected operation of the setf instruction (<4>).</p> </td></tr> </table> <p><u>Workaround</u></p> <p>(1) Do not use the sld instruction (e. g. by avoiding code optimization that makes use of sld).</p> <p>(2) If a code sequence as described above is used (a sld instruction following an instruction that can be executed in parallel), insert a nop instruction before the sld instruction.</p> <p>(3) If a code sequence as described above is used (a sld instruction following an instruction that can be executed in parallel), exchange the order of the previous two instructions as long as the program algorithm is not disturbed:</p>	<pre> <1> ld.w [r11], r10 : <2> mov r10, r28 <3> sld.w 0x28, r10 </pre>	<p>This situation occurs when the decode operation of the mov instruction (<2>) immediately before the sld instruction (<3>) and interrupt request servicing conflict before the execution of the special instruction ld (<1>) is complete.</p>	<pre> <1> ld.w [r11], r10 : <2> cmp imm5, r10 <3> sld.w 0x28, r10 <4> bz label </pre>	<p>This situation occurs when the decode operation of cmp (<2>) immediately before the sld instruction (<3>) and interrupt request servicing conflict before execution of the special instruction ld (<1>) is complete. As a result, the compare result of the cmp instruction becomes undefined, which may cause an unexpected operation of the branch instruction bz (<4>).</p>	<pre> <1> ld.w [r11], r10 : <2> add imm5, r10 <3> sld.w 0x28, r10 <4> setf c, r16 </pre>	<p>This situation occurs when the decode operation of the add instruction (<2>) immediately before the sld instruction (<3>) and interrupt request servicing conflict before execution of the special instruction ld (<1>) is complete. As a result, the result of the add instruction and the depending status flags become undefined, which may cause an unexpected operation of the setf instruction (<4>).</p>
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<pre> <1> ld.w [r11], r10 : <2> cmp imm5, r10 <3> sld.w 0x28, r10 <4> bz label </pre>	<p>This situation occurs when the decode operation of cmp (<2>) immediately before the sld instruction (<3>) and interrupt request servicing conflict before execution of the special instruction ld (<1>) is complete. As a result, the compare result of the cmp instruction becomes undefined, which may cause an unexpected operation of the branch instruction bz (<4>).</p>						
<pre> <1> ld.w [r11], r10 : <2> add imm5, r10 <3> sld.w 0x28, r10 <4> setf c, r16 </pre>	<p>This situation occurs when the decode operation of the add instruction (<2>) immediately before the sld instruction (<3>) and interrupt request servicing conflict before execution of the special instruction ld (<1>) is complete. As a result, the result of the add instruction and the depending status flags become undefined, which may cause an unexpected operation of the setf instruction (<4>).</p>						

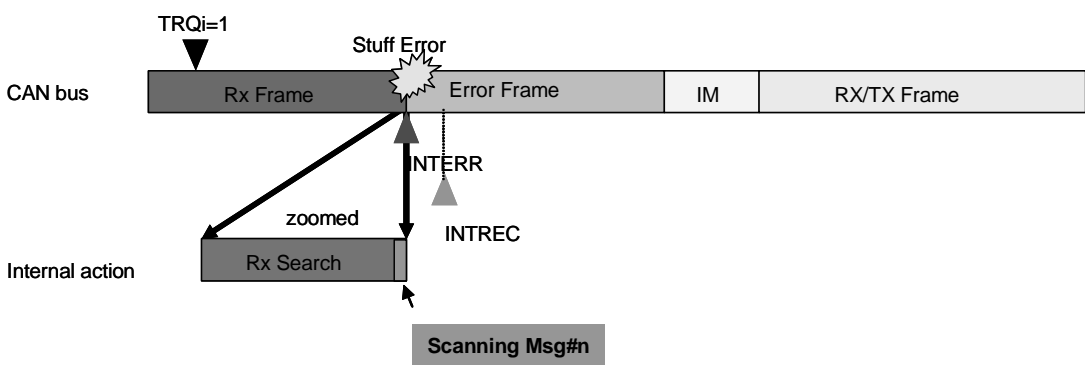
No. 13	SLD instruction precaution (Specification change notice) (Specification change notice)
	<p>(cont.)</p> <p>Example:</p> <p>1. (before implementing workaround)</p> <pre>ld.w [r11], r10 ... add r11, r12 mov r10, r28 sld.w 0x28, r10</pre> <p>2. (after implementing workaround)</p> <pre>ld.w [r11], r10 ... mov r10, r28 add r11, r12 sld.w 0x28, r10</pre> <p>(4) When assembler code is used: Avoid the critical code sequences as described above.</p>

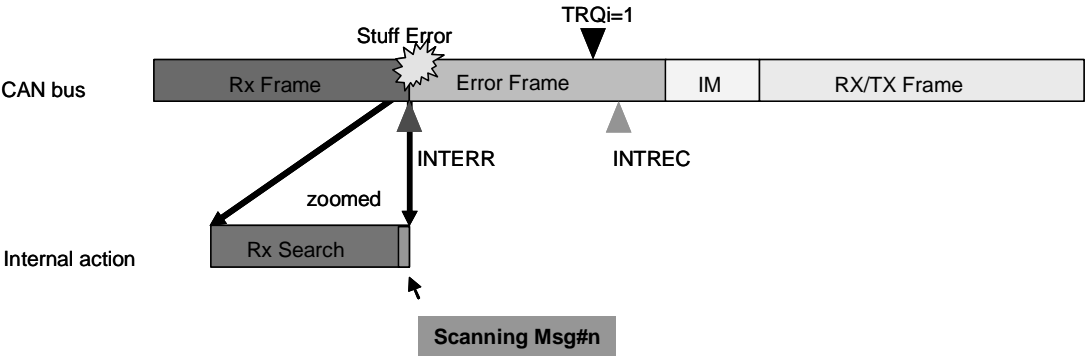
No. 14	aFCAN transmission / reception (Technical limitation)
	<p><u>Details</u></p> <p>The aFCAN macro will under certain timing conditions accompanied by a particular configuration of the message buffers and a specific operational usage not operate as expected. Different behaviors have to be considered. The description of the particular unexpected behavior is organized by the buffer Type: TX and RX, the frame format (extended or standard identifier), and the buffer number.</p> <p>TX-Buffer Behavior</p> <p>This section describes all unexpected behaviors linked to the configuration of transmit buffers. Configurations not listed are of no concern and can be used without restrictions.</p> <p>- Buffer #0 or Buffer #0 and #1 are configured as TX-buffer</p> <p>When using message buffer #0 as a TX-buffer, the message requested for transmission of this buffer may not be sent at the next possible timing. This behavior is caused when the internal scan for new transmission requests reached buffer #0 and at the same time the transmission of a previously sent message ends. Instead of sending the message object from buffer #0, the aFCAN attempts to send the contents of message buffer #1. In case the TRQ of buffer #1 is set, the message from buffer #1 is sent followed by the transmission of the message in buffer #0. This resembles an inner priority inversion. In case the TRQ of buffer #1 is not set or buffer #1 is not a TX-buffer, the contents of buffer #0 are sent whenever any of the other TRQ-bits in the aFCAN are set or cleared, or when a receive operation is started; i.e. when the next bus activity occurs. This behavior is valid for both frame types, extended or standard identifier format.</p> <p>- Buffer #1 - #31 are configured as TX-buffer</p> <p>There are two configurations that lead to the same, unexpected behavior. In the first configuration Buffer #1 through #31 are set up as normal TX-buffers, and in the second configuration buffer #0 through #7 are operated in ABT-mode (automatic block transmission) and the remaining buffers (#8 - 31) are configured as normal TX-buffers. The unexpected behavior occurs as well if only a subset of buffers are configured as normal TX-buffers. When using any message buffer #n in the range of buffer #1 through #31 with extended identifier, an inner priority inversion can be encountered. In that case the message from buffer #n+1 is sent in advance of the message in buffer #n even though the priority of the identifier in buffer #n is higher. This behavior is seen when the internal transmit search algorithms of the aFCAN processing buffer #n meets the start of a transmission, the end of a frame on the bus (RX or TX) that needs to end with an error frame, or the event of transmission request by the CPU or by the ABTmode. In case the ABT-mode is active, the unexpected behavior does only apply for messages in buffer #8 through #31. This behavior applies only if extended identifiers are in use. Applications exclusively using standard Identifiers do not suffer any limitation of this kind.</p>

No. 14	aFCAN transmission / reception (contd.)
	<p>RX-Buffer Behavior</p> <p>There are several configurations for the receive buffers that can cause different unexpected behaviors.</p> <ul style="list-style-type: none"> - Buffer #0 is configured as RX-buffer with EXT ID handling <p>There are two kinds of unexpected behavior in this configuration. The first one requires that a newly received message would normally be received in buffer #0. When in this case a transmission request (TRQ) by the host processor is submitted at the time where buffer #0 is scanned by the internal RX-Search algorithm for newly received messages, the storage of the message for buffer #0 will not occur. The message will be stored nowhere.</p> <p>In a second scenario, the newly received message would under normal conditions not be stored in buffer #0. When in this case a TRQ for buffer #i by the host processor is submitted, which again happens at the same time where the acceptance filtering for the newly received message is active, the message can be stored in buffer #0. The unexpected storage of the message in buffer #0 does only occur when the identifier bits ID15-0 of the received message coincidentally match the values in MCONF/MDLC of the TX-buffer #i. The behavior applies only for extended format frames. Standard format frames can be used without suffering this limitation.</p> <ul style="list-style-type: none"> - Buffer #1 - #31 are configured as RX-buffer with EXT ID handling <p>When a newly received message does not match the acceptance filter criteria for buffer #n (n = 1...31), the following timing can lead to an unexpected behavior. When additionally to that condition a transmission request is set for buffer #i at the time where buffer #n is scanned by the internal RX-Search algorithm, the storage of the message can be performed in buffer #n+1 although this buffer does not match with the acceptance filter criteria for the received message. The unexpected storage of the message in buffer #n+1 does only occur when the identifier bits ID15-0 of the received message coincidentally match the values in MCONF/MDLC of the Txbuffer #i. This unexpected behavior is also possible if only a subset of buffers in the range of #1 through #31 are configured as RX-buffers. The behavior applies only for extended format frames. Standard format frames can be used without suffering this limitation.</p> <p><u>Workaround</u></p> <p>The application needs to abstain from the usage of message buffer #0 or configure buffer #0 as a receive buffer with standard identifier handling.</p> <p>The ABT-mode is not affected from this limitation. Thus buffer #0 can be configured as a transmit buffer in that mode.</p> <p>Do not use extended identifiers for transmit and receive objects. If the application needs to process extended identifiers, the host processor must not issue a transmission request anytime the CAN-bus is busy. For this case the application can monitor the bus status and issue a transmission request right after the bus idle state was detected. Then the TRQ will be submitted in an uncritical point of time. This is even valid when the TRQ is submitted during the first 19 bits of the extended identifier.</p>

No. 15	TMPn / TMQn external event counter function (Direction of use)
	<p><u>Details</u></p> <p>When the external event counter mode is used and the compare register of timer TMP or TMQ is set to 0x0000 an interrupt occurs after the overflow of the timer.</p> <p><u>Workaround</u></p> <p>Avoid setting 0x0000 to the timer compare registers of TMP or TMQ.</p>

No. 16	TMPn / TMQn capture operation (Direction of use)
	<p><u>Details</u></p> <p>When the pulse width measurement mode or the free-running mode are used the initial value of the capture register is 0xFFFF when the count clock of the TMPn / TMQn is lower than the sampling clock of the capture trigger input after the timer is enabled (TPnCE = 1, TQnCE = 1).</p> <p><u>Workaround</u></p> <p>There is no workaround.</p>

No. 17	aFCAN: Rx limitation (Technical limitation)
	<p><u>Details</u></p> <p>RX Limitation</p> <p>The aFCAN macro may store an incoming message although this message was interrupted by a bus error frame. Thus, the incomplete reception causes that a message buffer is updated with old or incorrect data or that the message is even stored at an incorrect location.</p> <p>This unexpected behaviour affords that the bus error occurs in a certain relation to the currently present message on the bus. The critical time window starts at the sample point of the LSB of the DLC-field and lasts for the duration of an internal process in the aFCAN macro (RX-search). This time window usually lasts for a few bit times only. The actual length depends on the clock supply for the AFCAN, the CPU accesses during this period, the baud rate and the number of message buffers of the particular AFCAN macro.</p> <p>In this time window the RX-search evaluates the received identifier of the current message. When the bus error is detected within this window and when the RX-search has just scanned buffer #n for reception and found it is matching, the message will unexpectedly be treated as a received message. As the time window is limited as described above, only a stuff bit error occurring right in this window can cause this behaviour.</p> <p>There are two types of unexpected behavior for the RX limitation depending on the presence of pending transmission request (TRQ) for any other message buffer.</p> <p>1. Behaviour at pending TRQ (TRQi = 1)</p> <p>When the host processor has already submitted a transmit request (TRQ) for at least one buffer, the unexpected reception of the message will take place into the message buffer found by internal RX-search. This is the correct location to store the message i.e. the acceptance filter criteria are correctly fulfilled. However the data part will be updated with the contents of the shift register of the CAN protocol core. As this register is immediately stopped at detection of the bus error, the data provided to the message buffer can not be interpreted by the host processor.</p>  <p>Figure 1: Behavior at pending TRQ</p> <p>As during a regular reception, the RX-interrupt (if enabled) is generated and the application processes the message object.</p>

No. 17	aFCAN: Rx limitation (Technical limitation)
	<p>2. Behavior without pending TRQ (TRQi = 0)</p> <p>In case the host processor has not submitted a transmit request (TRQ) for any buffer before the detection of the bus error but submits TRQ = 1 after that point in time (see figure below) before the re-transmission of the message interrupted by the stuff bit error started, the unexpected reception of the message will take place.</p>  <p>Figure 2: Behavior without pending TRQ</p> <p>The unexpected storage of the message is issued in the particular message buffer that matches the acceptance filter criteria at the time where the bus error is detected (as described in 1.) or the message buffer #0 is overwritten independently of its configuration.</p> <p>Impact on application</p> <p>In typical applications the RX-limitation will lead to transiently incorrect data. In the vast majority of cases the message interrupted by a bus error is repeated by the transmitter right away. Then the application receives correct data shortly after the unexpected reception.</p> <p>In scenarios where the message buffer #0 is overwritten, the impact for the application depends on the usage of that buffer. If it is configured as a receive buffer, the application receives a message at an unexpected location and will interpret the data to belong to the identifier originally programmed for that buffer. The message buffer #0 needs to be re-configured in order to receive the originally intended message object again.</p> <p>In case of a transmit message buffer the unexpected storage may falsify a transmit object; i.e. when the unexpected behavior occurs after preparation of the message data but before the actual start of transmission. This scenario is even less likely than the scenario described in 1, which itself has a low probability. However the transmission of a falsified message can lead to repetitive transmission attempts when the original provider of that message (identifier) tries to send its message at the same time. Then the messages most likely will differ in their data part and a bit error is detected. This repetition resumes until one of the nodes enters error passive or bus off state. Then the situation is resolved as all pending TRQ are send with delay or are cancelled (in case of bus off state).</p>

No. 17	aFCAN: Rx limitation (Technical limitation)
	<p><u>Workaround</u></p> <p>NEC will update the affected products.</p> <p>NEC does not recommend a S/W workaround as first choice as it is fairly complex. On the one hand it is based on the control of submitting transmission requests only when the bus is idle. On the other hand a less complex algorithm can be used which does not prevent the unexpected reception but detects it safely and discards the unexpected reception in the CAN S/W driver. Any of these algorithms require that message buffer #0 is not used or that a 'dummy' TRQ in an unused buffer is set. This prevents behaviors as described in 2.</p>

(C) Valid Specification

Item	Date pulished	Document No.	Document Title
1	March, 2003	SUD-DT-02-0101-2-E	IE-703288-G1-EM1 (Preliminary User's Manual)
2	September, 2003	U16541EJ2V0UD00	V850ES/SG2 Hardware (Preliminary User's Manual)
3	November, 2003	U16603EJ2V0UD00	V850ES/SJ2 Hardware (Preliminary User's Manual)
4	April, 2004	U15943EJ3V0UM00	V850ES Architecture (User's Manual)

(D) Revision History

Item	Date pulished	Document No.	Comment
1	April 28, 2003	TPS-HE-B-2840	First release
2	August 11, 2003	TPS-HE-B-2841	Added item 4, Control Code 'D'
3	April 7, 2003	TPS-HE-B-2842	Added items 4 to 14, Control Code 'E'
4	August, 2004	TPS-HE-B-2843	Added items 15 to 17